

# SUPERMAN: Security Using Pre-Existing Routing for Mobile Ad hoc Networks

<sup>[1]</sup>R.Sangeetha, <sup>[2]</sup>R.Selvapriya, <sup>[3]</sup>M.Supraja, <sup>[4]</sup>D.Bhavanidevi  
<sup>[1][2][3]</sup>UG student, <sup>[4]</sup>Asst.Professor  
<sup>[1][2][3][4]</sup>Department of CSE, V.S.B Engineering College, Karur

**Abstract** – The adaptability and versatility of Mobile Ad hoc Networks (MANETs) have made them expanding famous in an extensive variety of utilization case. To protect these networks, security protocols have been developed to protect routing and application data. However, these protocols only protect routes or communication, not both. Both secure directing and correspondence security conventions must be executed to give full assurance. The use of communication security protocols originally developed for wire line and Wi-Fi networks can also place a heavy burden on the limited network resources of a MANET. To address these issues, a novel secure structure is proposed. The framework is designed to allow existing network and routing protocols to perform their functions, whilst providing node authentication, access control, and communication security mechanisms. This paper shows a novel security system for MANETs, SUPERMAN. Simulation results comparing SUPERMAN with IPSec, SAODV and SOLSR are provided to demonstrate the proposed frameworks suitability for wireless communication security.

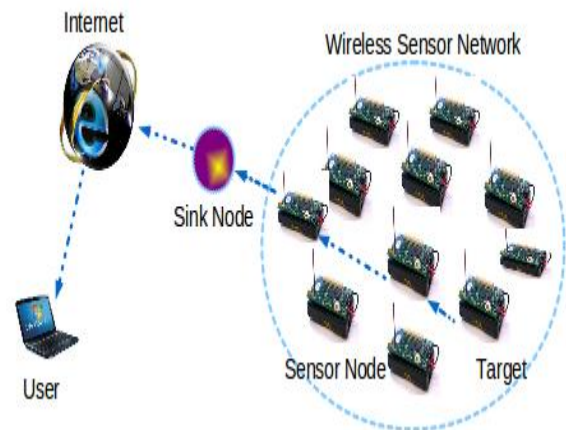
**Key words:** Data total system, Collusion assault, Wireless sensor systems, Iterative sifting calculation.

## 1. INTRODUCTION

A wireless sensor network (WSN) consists of a collection of these nodes that have the facility to sense, process data and communicate with each other via a wireless connection. Remote sensor systems (WSN's), the change in sensor innovation has made it conceivable to have little, low fueled detecting gadgets furnished with programmable process, numerous parameter detecting and remote message capacity. Likewise, the minimal effort makes it conceivable to have a system of hundreds or thousands of these sensors, along these lines upgrading the consistency and precision of information and the zone scope. Remote sensor systems offer data about confined structures, across the board natural changes, and so on. Remote sensor organize (WSN) is a system framework contained spatially disseminated gadgets utilizing remote sensor hubs to screen physical or natural circumstance, for example, sound, temperature, and movement.

Trust and reputation systems have a sign role in supporting operation of a wide range of distributed systems, from wireless sensor networks and e-commerce infrastructure to social networks, by providing an assessment of trustworthiness of participants in such distributed systems. A trustworthiness assessment at any given moment represents an aggregate of the behavior of the participants up to that moment and has to be robust in the presence of various types of faults and malicious behavior. There are a number of incentives for attackers to manipulate the trust and reputation scores of participants in a distributed system, and such

manipulation can severely impair the performance of such a system. The main target of malicious attackers is aggregation algorithms of trust and reputation systems.



**FIGURE 1: An operating system of a WSN**

A sensor network is designed to perform a set of high-level information processing tasks such as detection, track, or categorization. Measures of performance for these tasks are well defined, including discovery of false alarms or misses, classification errors, and track quality. As the computational power of very low power processors dramatically increases, mostly driven by demands of mobile computing, and as the cost of such technology drops, WSNs will be able to afford hardware which can implement more sophisticated data aggregation and trust assessment algorithms; an example is the recent

emergence of multi-core and multi-processor systems in sensor nodes. Iterative Filtering (PROPOSED) calculations are an alluring alternative for WSN's in light of the fact that they take care of the two issues - information total and information reliability evaluation - utilizing a solitary iterative technique. Such trustworthiness estimate of each sensor is based on the distance of the readings of such a sensor from the estimate of the correct values, obtained in the previous round of iteration by some form of aggregation of the readings of all sensors. Such aggregation is usually a weighted average; sensors whose readings sign proposed from such estimate are assigned less trustworthiness and consequently in the aggregation process in the present round of iteration their readings are given a lower weight.

## 2. RELATED WORKS

Maintaining a long network time with stringent energy constraints on tiny sensor nodes poses an extremely challenging task for wireless sensor networks. This paper provides a thorough analysis on a randomized algorithm that makes scheduling decisions without the help of geographic information. The analytical results precisely describe the relationship among achievable network coverage, energy saving, and node density. We likewise break down the execution of the randomized calculation with time asynchrony and propose a heuristic randomized booking plan to enhance the execution. The contribution of this paper is in four aspects. First, we build a mathematical model to analyze a purely randomized sensor node scheduling algorithm and to illustrate the relationship among achievable coverage quality, energy saving, and node density. Second, we prove that the purely randomized algorithm is resilient to time asynchrony proposed the network is sufficiently dense. Such feature is indispensable for a practical scheduling algorithm since precise time synchronization is very hard for large sensor networks. Our proof hence provides strong supporting evidence on the randomized scheduling algorithm for realistic applications. Third, we propose a heuristic strategy that enhances the achievable scope nature of the absolutely randomized calculation by equitably doling out neighboring sensor hubs into different checking sets. Finally, simulation study is performed to vary the correctness of the analytical results and to demonstrate the advantages of the proposed heuristic method. Propose a versatile Lightweight Deployment-Aware Scheduling calculation, which kills excess sensors without utilizing exact area data. Simulation study demonstrates that the LDAS algorithm

can reduce network energy consumption and provide desired QoS requirement effectively. Our analytical results will benefit the research in wireless sensor networks by providing simple formula to estimate sensor redundancy. They can be utilized in designing deployment-aware scheduling scheme to save energy consumption. Different sensor deployment strategies can cause very different network topology, and thus different degrees of sensor redundancy. The knowledge for sensor deployment, however, is usually available in advance. For instance, it is easy to know the number of sensors and how the sensors are dispatched for a particular application. Technical challenges in sensor network development include network discovery, control and routing, collaborative signal and information processing, tasking and querying, and security. The paper finishes up by showing some current research brings about sensor organize calculations, including limited calculations and coordinated diffusion, appropriated following in remote impromptu systems, and disseminated classification utilizing nearby specialists. Current and potential applications of sensor networks include: military sensing, physical security, air traffic control, traffic surveillance, video surveillance, industrial and manufacturing automation, distributed robotics, environment monitoring, and building and structures monitoring. The sensors in these applications might be little or huge, and the systems might be wired or remote. However, ubiquitous wireless networks of micro sensors probably offer the most potential in changing the world of sensing.

Simulated honey bee settlement (SUPERMAN) is the one which has been most generally considered on and connected to take care of this present reality issues, up until now. Day by day the number of researchers being interested in SUPERMAN algorithm increases rapidly. This work presents a comprehensive survey of the advances with SUPERMAN and its applications. It is hoped that this survey would be very beneficial for the researchers studying on SI, particularly SUPERMAN algorithm. The term swarm is used for an aggregation of animals like fishes, birds and insects such as ants, termites and bees performing collective behavior. The individual agents of these swarms behave without supervision and each of these agents has a stochastic behavior due to her perception in the neighborhood. SI is defined as the collective behavior of decentralized and self-organized swarms. Well known examples of which are bird flocks and the colony of social insects such as ants and bees. Results show artificial bee colony algorithm can be preferable in the dynamic deployment of wireless sensor

networks. This approach is based on Artificial Bee Colony (SUPERMAN) algorithm which is developed by modeling foraging behavior of honey bee swarms. It is known that the SUPERMAN algorithm works well for numerical optimization problems. The SUPERMAN calculation was first tried on powerful arrangement an utilizing parallel identification model of remote sensor systems comprising of every single versatile hub. Considering the good performance of the algorithm, use of the SUPERMAN algorithm will be a proper approach for the sensors in the network to obtain a good coverage in two dimensional spaces with stationary and mobile nodes. The execution of proposed approach is assessed in correlation with another swarm based system, Particle Swarm Optimization (PSO).

**2.1 PROBLEM DEFINITION**

The objective is

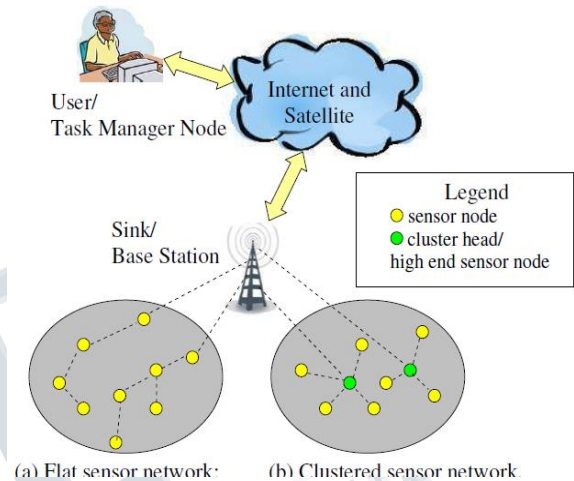
- 1) To deploy the sensor nodes such that the network time is maximum and
- 2) To schedule the sensor nodes so as to achieve the optimal network time.

**3. METHODOLOGY**

**3.1 NETWORK MODEL**

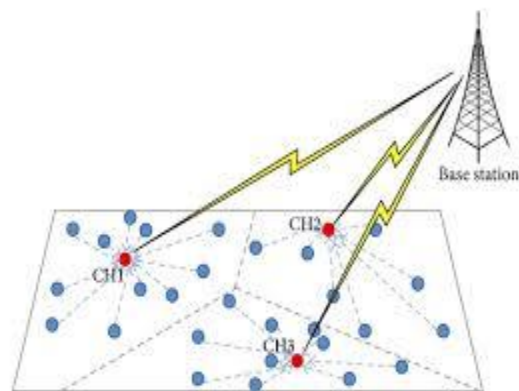
A WSN consists of small-sized sensor devices, which are equipped with limited battery power and are capable of wireless communications. When a WSN is deployed in a sensing field, these sensor nodes will be responsible for sensing abnormal events or for collecting the sensed data of the environment. In the case of a sensor node detecting an abnormal event or being set to periodically report the sensed data, it will send the message hop-by-hop to a special node, called a sink node. The sink node will then inform the supervisor through the Internet. Named data networking (NDN) resolves the traditional TCP/IP based Internet problems (i.e., location dependent, complex usage, scalability, poor resource utilization etc.) and is considered as an eligible candidate for futuristic Internet paradigm. In NDN-based mobile ad hoc networks (MANETs), the participating nodes are operated in highly dynamic and challenge able environment such as low battery power, channel fluctuations, intermittent connectivity and so on. Because of the communicate idea of the remote channel, the NDN-based MANETs feature serious issues (e.g., bundle crashes, flooding, information repetition, parcel retransmissions), which additionally debate the system execution. In this paper, to cope with these problems, we have proposed a novel protocol, named location-aware on-demand multipath caching and

forwarding (LOMCF) for NDN-based MANETs. Performance of the proposed protocol is evaluated by using simulator, called NDN SIM. Extensive experiments along their results show that proposed protocol performs better as comparing to the other recent proposed protocols in terms of content retrieval time, Interest retransmissions, and the total number of Interest packets injected as well as discarded in the network.



**FIGURE 2: Network model of a WSN**

The sensor nodes are divided into disjoint clusters, and every cluster has a cluster head which acts as an aggregator. Data are periodically together and aggregated by the aggregator.



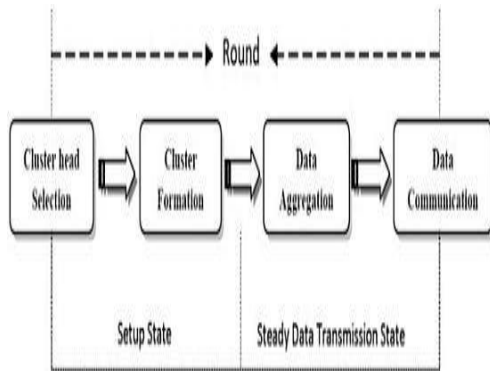
**FIGURE 3: Cluster head communication**

In this paper underestimate that the aggregator itself isn't bargained and focus on calculations which influence collection to secure when the individual sensor hubs may be traded off and may send false information to the aggregator. Assume that every data aggregator has

enough computational power to run an PROPOSED algorithm for data aggregation.

**3.2 FRAMEWORK OVERVIEW**

In order to recover the performance of PROPOSED algorithms against the aforementioned attack scenario, we provide a robust initial estimation of the trustworthiness of sensor nodes to be used in the first iteration of the PROPOSED algorithm.



**FIGURE 5: Framework overview of Data Aggregation Technique**

Most of the traditional statistical estimation methods for variance involve use of the sample mean. For this reason, proposing a robust variance estimation method in the case of skewed sample mean is an essential part of our methodology.

**3.3 ENHANCED ITERATIVE FILTERING ALGORITHM**

PROPOSED algorithm is robust against the simple outlier injection by the compromised nodes. An adversary employs three compromised

**OPTIMAL LOCATION DETERMINATION**

1) Sensor Deployment to Achieve 1-Coverage: Given a set of  $n$  targets  $T = \{T_1, T_2, \dots, T_n\}$  located in  $u \times v$  region and  $m$  sensor nodes  $S = \{S_1, S_2, \dots, S_m\}$ , place the nodes such that each target is monitored by at least one sensor node and the network time is maximum. The goal is to boost  $U$  to such an extent that each objective is observed by no less than one sensor hub.

2) Sensor Deployment to Achieve  $k$ -Coverage: Given an arrangement of  $n$  targets  $T = \{T_1, T_2, \dots, T_n\}$  situated in  $u \times v$  area and  $m$  sensor hubs  $S = \{S_1, S_2, \dots, S_m\}$ , put the hubs with the end goal that each objective is checked by in any event  $k$ -sensor hubs and to boost  $U$ .

3) Sensor Deployment to Achieve  $Q$ -Coverage: Given an arrangement of  $n$  targets  $T = \{T_1, T_2, \dots, T_n\}$  situated in  $u \times v$  area and  $m$  sensor hubs  $S = \{S_1, S_2, \dots, S_m\}$ , put the hubs with the end goal that each objective  $T_j, 1 \leq j \leq n$ , is secured by at any rate  $q_j$  sensor hubs and to expand  $U$ .

**3.4 PROCESS**

The nodes are initially deployed randomly. Based on the theoretical upper bound of network proposed time. We compute the optimal deployment locations using SUPERMAN algorithm. A heuristic is then used to schedule the sensor nodes such that the network proposed time is maximum.

**3.5 SCHEDULING ALGORITHM**

**3.5.1 SUPERMAN BASED SENSOR DEPLOYMENT**

SUPERMAN)Algorithm is an optimization algorithm based on the intelligent behavior of honey bee swarm. The colony of bees contains three groups: employed bees, onlookers and scouts. The employed bee takes a load of nectar from the source and returns to the hive and unloads the nectar to a food store. After unloading the food, the bee performs a special form of dance called waggle dance which contains information about the direction in which the food will be found, its distance from the hive and its quality rating. Since information about all the current rich sources is available to an onlooker on the dance floor, an onlooker bee probably could watch numerous dances and choose to employ itself at the most qualitative source. There is a more prominent likelihood of spectators picking more subjective sources since more data is flowing about the more subjective sources. Employed foragers share their information with a probability, which is proportional to the quality of the food source. Hence, the recruitment is proportional to quality of a food source. Exploitation is carried out by employed bees and onlookers, while exploration is carried out by scouts.

**3.5.2 PARTICLE SWARM OPTIMIZATION**

Particle Swarm Optimization (PSO) consists of a swarm of particles moving in a search space of possible solutions for a problem. Every particle has a position vector representing a candidate solution to the problem and a velocity vector. Moreover, each particle contains a small memory that stores its own best position seen so far and a global best position obtained through communication with its neighbor particles. It consists of a swarm of  $w$  candidate solutions called particles, which explore an  $n$ -dimensional hyperspace in search of the global solution ( $n$  represents the number of optimal parameters to be

determined). A particle  $p$  occupies position  $x_{pd}$  and velocity  $v_{pd}$  in the  $d$ th dimension of the hyperspace,  $1 \leq p \leq w$  and  $1 \leq d \leq nd$ . In the global-best version of PSO, the position where the particle  $p$  has its best cost is stored as  $(pbestpd)$ . Besides,  $gbestd$ , the position of the best particle. In each iteration  $tr$ , velocity  $v$  and position  $x$  are updated using (5) and (6). The update process is iteratively repeated until either an acceptable best is achieved or a fixed number of iterations is reached.

### 3.5.3 SUPERMAN

The flexibility and mobility of Mobile Ad hoc Networks (MANETs) have made them increasing popular in a wide range of use cases. To ensure these systems, security conventions have been created to secure steering and application information. However, these protocols only protect routes or communication, not both. Both secure routing and communication security protocols must be implemented to provide full protection. The use of communication security conventions initially created for wireline and WiFi systems can likewise put an overwhelming weight on the constrained system assets of a MANET. To address these issues, a novel secure framework (SUPERMAN) is proposed. The structure is intended to enable existing system and steering conventions to play out their capacities, while giving hub verification, get to control, and correspondence security instruments. This paper presents a novel security framework for MANETs, SUPERMAN. Amusement happens differentiating SUPERMAN and IPSec, SAODV and SOLSR are given to demonstrate the proposed structures sensibility for remote correspondence security.

### 3.6 HEURISTIC FOR SENSOR SCHEDULING

To achieve this, we propose a weight-based method for determining the cover sets

- 1) Weight assignment
- 2) Cover formation
- 3) Cover optimization
- 4) Cover activation and Energy reduction.

## 4. EXPERIMENTAL RESULTS

The objective of our experiments is to evaluate the robustness and efficiency of our approach for estimating the true values of signal based on the sensor readings in the presence of faults and collusion attacks. For each experiment, we evaluate the accuracy based on Root Mean Squared error (RMS error) metric and efficiency based on the number of iterations needed for convergence of proposed algorithms.

The main shortcoming of the proposed algorithms in the proposed attack scenario is that they quickly converge to the sample mean in the presence of the attack scenario. Keeping in mind the end goal to explore the inadequacy, we led a test by expanding the sensor fluctuations and additionally the quantity of colluders.

In this experiment, quant proposed IED the number of iterations for the proposed algorithm with reciprocal discriminate function. The results obtain from this experiment show that the original version of the proposed algorithm quickly converges to the skewed values provided by one of the attackers, while starting with an initial reputation provided by our approach, the algorithms require around 29 iterations, and, instead of converging to the skewed value provided by one of the attackers, it provide a reasonable accuracy. The results of this experiment show that the proposed initial reputation for the PROPOSED algorithm improve the efficiency of the algorithm in terms of the number of iterations until the process has converged. In other words, by providing this initial reputation, the number of iterations for proposed algorithm decreases approximately 9% for reciprocal and around 8% for affine discriminant functions in both biased and unbiased conditions. This can be explained by the fact that the new initial reputation is close to the true value of signal and the proposed algorithm needs fewer iterations to reach its stationary point.

## 5. CONCLUSION

Mobile Ad-hoc Networks (MANETs) are a cluster of self-organizing and self-governing wireless nodes without any backbone infrastructure and centralized administration. The different hubs in MANET move arbitrarily, and this hub portability may posture challenges on the execution of directing conventions. In this paper, an Intra and intergroup performance review of various MANET routing protocols are performed under varying speed of nodes. The routing protocols included in this study are reactive, proactive, and hybrid protocols. This performance review is done using the Java simulator and random way point model. The routing protocols performance is assessed through standard performance measure metrics including packet delivery ratio, throughput, routing overhead and end to end delivery with varying speed of nodes. The reproductions result demonstrates that there is no critical effect of fluctuating rate of hubs on standard execution assessment measurements.

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