

An Efficient and Secure Anti-Collusion Data Sharing Scheme for Dynamic Groups in the Cloud

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Abstract: — In cloud giving security, ensures for the sharing information record. Sadly, in view of the continuous variation of the participation, allocation files while giving protection careful is so far a difficult issue, specially for an untrusted cloud on account of the arrangement assault. In this exploration work, we propose an ensured information sharing course of action for dynamic individuals. Right off the bat, we propose a protected course for key appointment with no secured correspondence channels, and the clients can safely get their private keys from social event boss. Moreover, our course of action can complete fine-grained find the opportunity to control, any client in the party can utilize the source in the cloud and denied clients can't get to the cloud again after they are repudiated. Thirdly, we can shield the plan from intrigue strike, which proposes that deny clients can't get the exceptional information record paying little regard to the probability that they outline with the untrusted cloud. This course of action can accomplish fine practicality, which gathers past clients require not to revive their private keys for the circumstance either another client appreciates the social event or a client is denied from the party.

Index Terms—cloud, key appropriation, fine-grained.

1. INTRODUCTION

Distributed computing, through the features of inherent information distribution and low maintenance, gives an unrivalled consumption of benefits. In distributed calculating, cloud expert organizations proposal a consideration of unending storage room for clients to have data. It can empower clients to diminish their cash related above of information administrations by affecting the close-by organizations group into cloud servers.

However, safekeeping alarms turn into the essential limitation as we now outsource the size of data, which is potentially delicate, to cloud sources. To safeguard data security, a typical method is to encode data files before the clients assignment the encoded data into the cloud. Unfortunately, it is harden to plan a protected and proficient data allocation schemes, exclusively for dynamic gatherings in the cloud.

Regardless, inquire about outcomes demonstrate that security concerns particularly data security and privacy protection issues, will remain the real reason of complaint for decision makers to adopt cloud computing.

Application programming databases are moved now towards cloud computing storage where clients may not feel that it is sufficiently dependable as indicated by these components: Trust administration, Security provider, Privacy assurance, Ownership Data area and Relocation, Data integrity, Data recuperation, Performance and accessibility, Data Backup, Data portability and change. Securing the data stored in the cloud is very important to organizations and enterprises before moving their urgent data from their in-premises facilitating to the cloud.

Our research is about another procedure for encryption data and storing them into a cloud storage. The commitment is about making a more secure data storage into the cloud

II. RELATED WORKS AND PRELIMINARIES

2.1. Secure Data Sharing in the Cloud

In this area, we talk about the developing requirement for information sharing and the advantages of information sharing by means of the Cloud. We incline the necessities of information partaking in the Cloud taken after by the customary way to deal with sharing information by means of the Cloud and why this isn't compelling. We likewise examine the key administration issue and audit various works that address this issue. We at that point audit late mechanism that mean to give private and protected information partaking in the Cloud and examine the most recent systems recycled to accomplish this.

2.2. Importance of Data Sharing

Information sharing is winding up noticeably progressively essential for some clients and once in a while a urgent prerequisite, particularly for organizations and associations expecting to pick up benefit. Generally, many individuals saw the PC as "unoriginal goliaths" who undermined to slice employments of many individuals through robotization. Notwithstanding, as of late, it has been invited by countless as it has turned out to be fundamentally social. It is accordingly not shocking that an ever increasing number of individuals are requesting

information sharing capacity on their telephones, PCs and even as of late Smart TVs.

With the headways in Cloud figuring, there is presently a developing spotlight on actualizing information sharing abilities in the Cloud.

With the capacity to share information through the Cloud, the quantity of advantages increments multi overlap. As organizations and associations are currently outsourcing information and processes to the Cloud, they event advance with the capacity to share information between different organizations and associations. Representatives additionally advantage as they can impart work and team up to different workers and can likewise keep functioning at home or whatever other residence, for example, the library. They don't have to stress over losing act as it is dependably in the Cloud. With social clients, the capacity to share records, containing archives, photographs and recordings with different clients gives incredible advantage to them.

2.3 Requirements of Data Sharing in the Cloud

To empower information partaking in the Cloud, it is basic that lone approved clients can access information put away in the Cloud. We abridge the perfect prerequisites of information partaking in the Cloud underneath.

- The information administrator should to have the capacity to indicate a gathering of clients that are permitted to see his/her information.
- Any individual from the gathering should access the information whenever without the information proprietor's intercession.
- No other client, other than the information proprietor and the individuals from the gathering, should access the information, including the Cloud Service Provider.
- The information administrator should to have the capacity to renounce access to information for any individual from the gathering.
- The information administrator should to have the capacity to add individuals to the gathering.
- No individual from the gathering ought to be permitted to repudiate privileges of different individuals from the gathering or join new clients to the gathering.
- The information owner should to have the capacity to indicate who has read/write

authorizations on the information owner's records.

We now appearance at the glance at the protection and security prerequisite of information partaking in the Cloud. Accomplishing these necessities in the Cloud design can go far to pulling in vast quantities of clients to receiving and grasping Cloud innovation.

- **Information Confidentiality:** Unauthorized clients (counting the Cloud), should not have the capacity to get to information at any given time. Information ought to stay classified in travel, very still and on reinforcement media. Just approved clients ought to have the capacity to access information.
- **Client repudiation:** When a client is renounced get to rights to information, that client should not have the capacity to access the information at any given time. In a perfect world, client repudiation should not influence other approved clients in the gathering for productivity purposes.
- **Adaptable and Effective:** Since the quantity of Cloud clients has a tendency to be amazingly extensive and on occasion flighty as clients join and abandon, it is basic that the framework keep up productivity and also be versatility
- **Collusion between substances:** When considering information sharing philosophies in the Cloud, it is key that notwithstanding when certain elements plot, they would in any case not have the capacity to get to any of the information proprietor's consent. Prior works of writing on information sharing did not deliberate this issue, however agreement between elements can never be composed off as an unconvincing occasion.

2.4 Collusion Attack:

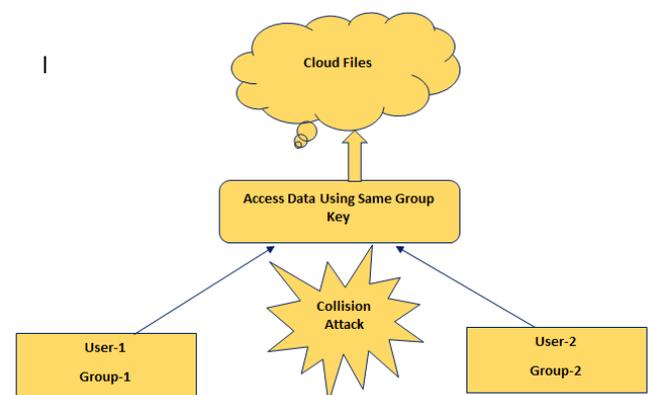


Fig 2.3.1: Collusion attack

From the above figure, we can notice that user-1 from the group-1 and user-2 from the group-2 both are accessing data sharing files by using the same group key. At that time collusion attack between two groups occurs. And this attack will overcome by giving separate group keys for different groups.

Preliminaries

Bilinear Maps

Give G_1 and G_2 a chance to be added substance cyclic gatherings of a similar prime request q .

Let $e : G_1 \times G_1 \rightarrow G_2$ indicate a bilinear guide built with the accompanying properties:

- 1) Bilinear: For all $a, b \in \mathbb{Z}_q^*$ and $P, Q \in G_1, e(aP, bQ) = e(P, Q)^{ab}$
- 2) Nondegenerate: There exists a point Q with the end goal that $e(Q, Q) \neq 1$.
- 3) Computable: There is an efficient algorithm to figure $e(P, Q)$ for any $P, Q \in G_1$

Key Generation

Step 1: Choose two particular prime numbers p and q .

Step 2: Find n to such an extent that $n = pq$.
 n will be utilized as the modulus for both the public and private keys.

Step 3: Find the totient of n , $\phi(n)$

$$\phi(n) = (p-1)(q-1).$$

Step 4: Choose an e with the end goal that $1 < e < \phi(n)$, and such that e and $\phi(n)$ share no divisors other than 1 (e and $\phi(n)$ are generally prime).

Step 5: Determine d (utilizing secluded math) which fulfills the congruence connection.

$$de = 1 \pmod{\phi(n)}.$$

As it were, pick d with the end goal that $de - 1$ can be equitably separated by $(p-1)(q-1)$, the totient, or $\phi(n)$.

This is regularly enlisted using the Extended Euclidean Algorithm, since e and $\phi(n)$ are modestly prime and d is to be the detached multiplicative turn around of e .

d is saved as the private key type.

Public key has modulus n and people in general (or encryption) sort e . The private key has modulus n and the private (or unscrambling) case d , which is preserved secret.

Encryption

Step 1: Person A communicates his/her public key (modulus n and example e) to Person B, keeping his/her private key mystery.

Step 2: At the point when Person B wishes to send the message "M" to Person A, he initially changes over M to a whole number with the end goal that $0 < m < n$ by utilizing settled upon reversible convention known as a padding plan.

Step 3: Person B calculates, with Person A's public key info, the cipher text c equivalent to $c = me \pmod{n}$.

Step 4: Person B now send message "M" in Cipher text, or c , to Person A.

Decryption

Step 1: Person A recovers m from c by using his/her Private key exponent, d , by the computation $m = cd \pmod{n}$.

Step 2: Given m , Person A can recover the original message "M" by withdrawing the padding scheme.

This method works since,

$$c = me \pmod{n}$$

$$cd = (me)d \pmod{n}$$

$$cd = mde \pmod{n}$$

By symmetry property of mods we have that $mde = m \pmod{n}$

Since $de = 1 + k\phi(n)$, we can write

$$mde = m(1 + k\phi(n)) \pmod{n}$$

$$mde = m(mk)\phi(n) \pmod{n}$$

$$mde = m \pmod{n}$$

From Euler's Theorem and the Chinese Remainder

Theorem, we can show that this is correct for all m and the original message

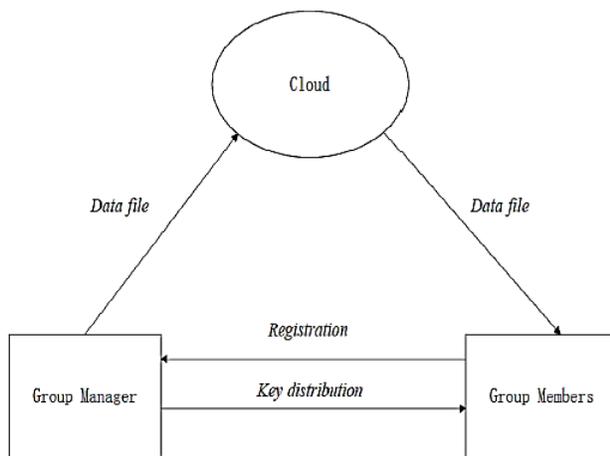
$$cd = m \pmod{n}, \text{ is obtained.}$$

III. EXISTING SYSTEM

- In this paper, we suggest a protected information sharing plan, which can accomplish secure key dissemination and information sharing for dynamic gathering.
- We give a safe approach to key appropriation with no safe correspondence channels. The customers can securely gain their private keys from bundle boss with no Certificate Authorities due to the check for the overall public key of the customer.
- Our plan can achieve fine-grained get to control, with the assist of the collecting client list, any client in the gathering can utilize the source in the cloud and revoked clients can't get to the cloud over after they are revoked.

- We propose a sheltered data sharing arrangement which can be protected from plot attack. The repudiated customers can not have the ability to get the primary data records once they are disavowed paying little mind to the likelihood that they imagine with the untrusted cloud. Our arrangement can achieve secure customer denial with the assistance of polynomial function.
- Our scheme is able to maintenance dynamic groups efficiently, when a new user joins in the group or a user is revoked from the group, the private keys of the other customers do not need to be recomputed and updated.
- We provide security analysis to prove the security of our scheme.

VI. SYSTEM ARCHITECTURE & DESIGN



The main aim is to provide privacy and security of group sharing data in public cloud computing.. The scope is to combine proxy signature, enhanced TGDH (Tree-Based Group Diffie-Hellman) and proxy re-encryption together into a protocol to effectively grant the privilege of group management and negotiate and update the group key pairs.

4.2. Scheme description

System Initialization

The group manager assumes responsibility of this operation. He creates a bilinear guide group framework $S=(q,G1,G2,e(...))$, then chooses two arbitrary components $P,G \in G1$ and a number $\gamma \in \mathbb{Z}^*$, at that point processes $W=\gamma.P, Y=\gamma.G$ and $Z=e(G,P)$. Finally, the group manager distributes the factors $(S,P,W,Y,Z,f1,Enc())$, where f is

hash function: $\{0,1\} \rightarrow \mathbb{Z}_q^*, f1$ is hash function: $\{0,1\}^* \rightarrow G1$, and $Enc()$ is a symmetric encryption calculation. Moreover, the gathering supervisor will keep the parameters (γ,G) as the secret key.

Registration for Existing User

This operation is performed by client(user), the group manager and the cloud. As a matter of first importance, the client sends $Idi, pk, v1$ as a demand to the group manager, where Idi is the identity of the client, pk is the public key utilized as a asymmetric encryption algorithm, for example, ECC, air conditioning is the record client used to pay for the registration, which is identified with the personality of the client, and $v1 \in \mathbb{Z}_q^*$ is an arbitrary number chosen by the client. The group manager compares the received Idi message and the personality Idi processed by decoding $AENCsk(Idi,vi,ac)$. Also, the group manager verifies if the decrypted number $v1$ is equivalent to the random number $v1$ in first step. After successful verifications, the group administrator produces the message KEY as takes after when the Idi message coordinates the individuality. At long last, the client decodes the message $AENCpk(KEY,v2)$ by his private key in ECC and afterward he can acquire his private key (xi,Ai,Bi) . After effective registration, the client turns into a group member.

File Upload

The group manager encodes cipher text $CE=\{C1,C2,C\}kr$ with the re-encryption key and sends $\{DF=(Idgroup,IDdata,CE,EK,tdata),\sigma DFg\}$ to the cloud, where $tdata$ is the time that the information file is transferred and $\sigma DF=\gamma f1(DF)$ is the mark of the group manager for the information file. What's more, the group manager likewise sends the information rundown to the cloud keeping in mind the end goal to give the clients a chance to check the freshness of the information file. the group manager includes $(Iddata,tdata)$ and the present time tDL to the information list DL . To guarantee that the two customers can secure the latest adjustment of data file and the cloud can invigorate the data file, the gathering chief revives the data list standard. Finally, the group executive includes his signature $sig(DL)=\gamma f1(DL)$ to the information list and sends the information rundown to the cloud for capacity. At long last, on getting the message, the cloud verifies the identity of the gathering director by checking the condition $e(W,f1(DF))=e(p,\sigma DF)$ and stores the message after effective verification.

User Revocation

At the point when a client i with identity ID_i is revoked, the group manager plays out the following operations:

- 1) Customer i should be removed from the gathering customer list in the area storage space and refreshing the gathering customer list which is secured in the cloud.
- 2) Testing the new group client list, assume that there are m appropriate group individuals in the rundown. As per the rundown, group manager at that point builds the new polynomial function $f_p(x) = \prod_{j=1}^m (x - V_j)$ and the new exponential function $\{W_0, \dots, W_{m-1}\} = \{G^a, \dots, G^{a(m-1)}\}$ where $G \in G_1$.
- 3) Selecting another arbitrary re-encryption key K^r and developing $EK = \{k^r, W_0, \dots, W_{m-1}\}$.
- 4) Computing cipher-text $CE = \{C_1, C_2, \dots, C_m\} K^r$ with the new re-encryption key K^r .
- 5) Validation his signature $\sigma(DF)$ to the modified message $DF = (ID_{group}, ID_{data}, CE, EK, t^1_{data})$, where t^1_{data} is the time stamp.

Registration for New User

The enrollment of another client with its identity ID_{m+1} , the group manager plays out an indistinguishable operation from enlistment for existing client. Likewise, the group manager refreshes the information files put away in the cloud. Above all else, the group manager checks whatever remains of the legal clients, at that point he builds the new polynomial function $f_p(x) = \prod_{j=1}^m (x - V_j)$ and the new exponential function W_0 . Finally, the group manager refreshes all the time cast t^1_{data} of the information files in the group for the information list. At that point the group manager directs the new information rundown to the cloud for storage.

File download

Having gotten the communication sent by the cloud, the aggregate part confirms the legitimacy of the information file and the rundown by testing the condition $(W, f_l(DF)) = e(P, \sigma(DF))$ and $(W, f_l(DL)) = e(P, \text{sig}(DL))$. then the gathering part orders if the time stamp put away in the DF and the information incline is same. At last the gathering part begins to decode the information file next fruitful confirmation.

V. HELPFUL HINTS

1. system initialization

In this step, the group manager assumes responsibility of this operation. He produces a bilinear guide bunch framework. at that point he will preserve the factors as the master key.

2. Registration for Existing user

This operation is performed by user, group manager, the cloud. user sends the registration request to the group manager. Then he can obtain his private key. After successful registration, the user becomes a group member.

3. File upload

In this operation, group member encrypts data to the group manager. Then he checks the data and upload the file to the cloud.

4. User revocation

This procedure is accomplished by the gathering administrator and the cloud. Here cloud substitutes the old information file with the new information file.

5. Registration for new user

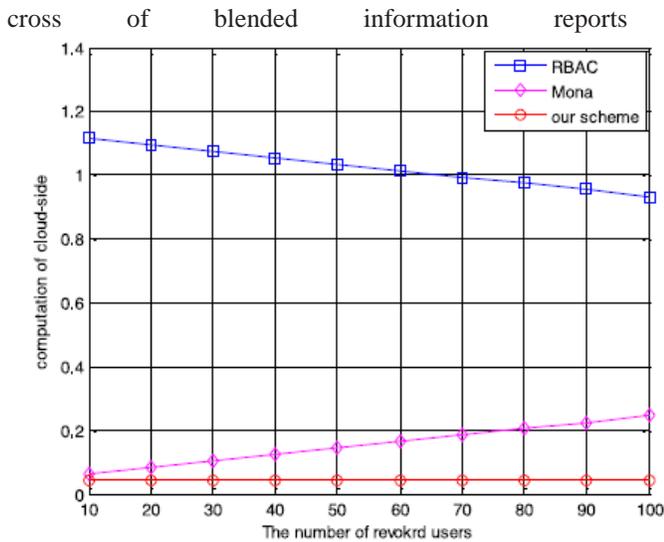
The group manager plays out an indistinguishable operations from enrollment for existing client, at that point he sends the original data once-over to the cloud for limit.

6. File download

This operation is performed by the group member. He can decrypt or download the data files from the cloud.

VI. RESULTS AND DISCUSSION

The calculation cost of the cloud for report download actions with the level of 100 Mbytes are addressed in Fig. 7.1. Like the operation of archive exchange, the count rate of the cloud is generally managed by the foreswearing check process. Along these lines, the cost increases with the amount of denied customers. In any case, in our arrangement, the cloud just basically checks the stamp. Along these lines, the number cost of the cloud for record download is unimportant to the measure of the repudiated clients. The illumination behind the high calculation cost of the cloud in RBAC plot is that the cloud plays out some tally operations to help the client to interpret information records. What's more, it can be seen that in these plans, the calculation cost is free with the measure of the chronicle, since both the stamp in Mona and the encoded message in our course of action are silly to the level of the asked for record and the operations of cloud for translating in RBAC plot is besides unessential to the



Downloading of 100MB file

Fig 7.1. Examination on calculation rate of the cloud for document download among RBAC, Mona and our plan.

VII. CONCLUSION

This plan outlines a secure anti-collusion data distribution plan for dynamic groups in the cloud. In this plan, the clients can firmly acquire their private keys from group manager with no Certificate Authorities and secure correspondence channels. Likewise, this plan can support dynamic groups productively, when another client contributes in the group or a client is revoked from the group, the private keys of substitute clients don't should be recomputed and relaxed. Also, this arrangement can achieve secure customer disavowal, the repudiated customers can not have the ability to get the principal information records once they are revoked paying little mind to the likelihood that they plot with the untrusted cloud.

In this research work, we have surveyed writing on approaches to give a secure situation where a data owner can impart data to individuals from his group while preventing any untouchables from increasing any data access if there should be an occurrence of malicious exercises, for example, data misfortune and robbery. Be that as it may, all through this work we expect that individuals from the group won't complete malicious exercises on the data owner's data.

VIII. FUTURE ENHANCEMENT

Auditing and Accountability in the Cloud is a potential for future research with regards to data sharing in the Cloud. Numerous clients specifically associations and undertakings pick up the advantage from data sharing in the Cloud. In any case, there is dependably a probable possibility that individuals from the group can complete unlawful operations on the data, for example, making illicit duplicates and circulating duplicates to companions, overall population, and so forth keeping in mind the end goal to benefit. A future research bearing is discover courses for a data owner to consider responsible any part that completes malicious exercises on their data.

Another research direction would be to give the data owner physical access control over the data. Rather than responsibility, the data owner can make an arrangement of get to control leads on his information and send the information alongside the get to control policy. In this way, any part with access to the information can just utilize the information in such a route that abides by the get to control policy. In the event that a member attempts to make illicit duplicates of the information, the get to control policy should "bolt" the information to keep the part from doing as such. Likewise, since information put away in the Cloud are normally put away and reproduced in various land areas around the globe, it is pivotal that the legal jurisdictions are respected and taken after. A potential research bearing is discover approaches to store and process information in a way that does not rupture the protection and security laws of the region.

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